

Example of Metropolis-Hastings (MH) Algorithm: Beta Distribution $B(\alpha, \beta)$:

$$f(x) = \begin{cases} \frac{1}{B(\alpha, \beta)} x^{\alpha-1} (1-x)^{\beta-1}, & \text{for } 0 < x < 1, \\ 0, & \text{otherwise,} \end{cases}$$

for $\alpha > 0$ and $\beta > 0$.

$f_*(x)$ is taken as the uniform distribution between zero and one:

$$f_*(x) = \begin{cases} 1, & \text{for } 0 < x < 1, \\ 0, & \text{otherwise,} \end{cases}$$

$$q(x) = \frac{f(x)}{f_*(x)} = \frac{1}{B(\alpha, \beta)} x^{\alpha-1} (1-x)^{\beta-1}$$

$\omega(x_{i-1}, x^*)$ is given by:

$$\omega(x_{i-1}, x^*) = \frac{q(x^*)}{q(x_{i-1})} = \frac{x^{*\alpha-1} (1-x^*)^{\beta-1}}{x_{i-1}^{\alpha-1} (1-x_{i-1})^{\beta-1}}$$

From computational viewpoint, $\omega(x_{i-1}, x^*)$ is computed as:

$$\omega(x_{i-1}, x^*) = \exp\left((\alpha - 1)\log(x^*) + (\beta - 1)\log(1 - x^*) - (\alpha - 1)\log(x_{i-1}) - (\beta - 1)\log(1 - x_{i-1})\right)$$

is recommended.

————— B(a,b) Distribution —————

```
1: #include <math.h>
2: #include <stdio.h>
3:
4:  int ix=1,iy=1;
5:
6: void main(){
7:
8:  float  a,b;
9:  long int i,j,n,m;
10:  double urnd(void);
11:  double x,x0,u,w;
```

```

12: double x1=0.0,x2=0.0;
13:
14: for(i=1;i<=10000;i++) urnd();
15:
16: scanf("%f%f%ld%ld",&a,&b,&m,&n);
17:
18: x=0.5;
19: for(i=-m+1;i<=n;i++){
20:     x0=urnd();
21:     w=exp( (a-1.)*log(x0)+(b-1.)*log(1.-x0)
22:           -(a-1.)*log(x )-(b-1.)*log(1.-x ) );
23:     u=urnd();
24:     if( u<w ) x=x0;
25:     if( i >= 1 ){
26:         x1+=x/((double)n);
27:         x2+=x*x/((double)n);
28:     }
29: }
30:
31: printf("# of Burn-in      = %10ld\n",m);

```

```

32: printf("# of Random Draws = %10ld\n",n);
33: printf("Parameters = (%7.1f,%7.1f)\n",a,b);
34: printf("Mean      = %10.5lf, which should be close to %10.5f\n"
35:        ,x1,a/(a+b));
36: printf("Variance = %10.5lf, which should be close to %10.5f\n"
37:        ,x2-x1*x1,a*b/((a+b)*(a+b)*(a+b+1.)));
38:
39: }
40: /* ----- */
41: double urnd(void)
42: {
43:     int    kx,ky;
44:     double rn;
45: /*
46:     Input:
47:         ix, iy:  Seeds
48:     Output:
49:         rn: Uniform Random Draw U(0,1)
50: */
51:     kx=ix/53668;

```

```

52:   ix=40014*(ix-kx*53668)-kx*12211;
53:
54:   ky=iy/52774;
55:   iy=40692*(iy-ky*52774)-ky*3791;
56:
57:   rn=(float)(ix-iy)/2147483563.;
58:   rn-=(int)rn;
59:   if( rn<0.) rn++;
60:
61:   return rn;
62: }

```

In Lines 18, an appropriate value is given to **x** as an initial value, i.e., x_{-M+1} , where M is **m**.

In Lines 19 – 29, **i** moves from **-m+1** to **n**.

The first **m** random draws are discarded, and the second **n** ones are utilized for further analysis such as Lines 25 – 28.

w in Line 21 corresponds to the acceptance probability $\omega(x_{i-1}, x^*)$, where x_{i-1} and x^* are given by x and x_0 in Line 20, respectively

In Line 24, x is updated to x_0 with probability w .

Remark: Theoretically, the acceptance probability w should be less than or equal to one, i.e., $w = \min(w, 1.)$ have to be included between Lines 22 – 23.

However, as in Line 24, if w is greater than one, x is always updated to x_0 .

In this case, even if the maximum value of w is restricted to one, x is always updated to x_0 .

Therefore, the sentence $w = \min(w, 1.)$ is redundant.

2.8.4 Sort Algorithm

In Bayesian analysis, sorting a sequence of n random draws x_1, x_2, \dots, x_n , quantiles are obtained.

Therefore, a sort program is very important.

There are a lot of sort algorithms (for example,

see <https://shinoarchive.com/contents/1916/>).

Insertion Sort: The simplest algorithm, called the insertion sort (挿入ソート), is:

————— Very Simple Sort Algorithm —————

```
1: #include<stdio.h>
2: #include<math.h>
3: #include<time.h>
4:
```

```
5:  int ix=1,iy=1;
6:  float y[100001];
7:
8:  void main()
9:  {
10:   int i,n;
11:   int i025,i050,i500,i950,i975;
12:   float x[100001];
13:   float urnd(void);
14:   void sort(float x[],int n);
15:   clock_t t0,t1;
16:   double dt;
17:
18:   for(i=1;i<=10000;i++) urnd();
19:
20:   scanf("%d",&n);
21:
22:   for(i=1;i<=n;i++){
23:     x[i]=urnd();
24:     y[i]=x[i];
```

```

25:  }
26:
27:  t0=clock();
28:  sort(x,n);
29:  t1=clock();
30:  dt=(t1-t0)/((double)CLOCKS_PER_SEC);
31: /*
32:  for(i=1;i<=n;i++) printf("%10d %10.8f %10.8f\n",i,x[i],y[i]);
33: */
34:  printf("Computational Time = %10.2lf\n",dt);
35:  i025=(int)( 0.025*(float)n );
36:  i050=(int)( 0.050*(float)n );
37:  i500=(int)( 0.500*(float)n );
38:  i950=(int)( 0.950*(float)n );
39:  i975=(int)( 0.975*(float)n );
40:  printf(" 2.5 percent point = %10.8f\n", (y[i025]+y[i025+1])/2.0 );
41:  printf(" 5 percent point = %10.8f\n", (y[i050]+y[i050+1])/2.0 );
42:  printf("50 percent point = %10.8f\n", (y[i500]+y[i500+1])/2.0 );
43:  printf("95 percent point = %10.8f\n", (y[i950]+y[i950+1])/2.0 );
44:  printf("97.5 percent point = %10.8f\n", (y[i975]+y[i975+1])/2.0 );

```

```

45: }
46: /* ===== */
47: float urnd(void)
48: {
49:     int    kx,ky;
50:     float  rn;
51: /*
52:     Input:
53:     ix, iy: Seeds
54:     Output:
55:     rn: Uniform Random Draw U(0,1)
56: */
57:     kx=ix/53668;
58:     ix=40014*(ix-kx*53668)-kx*12211;
59:
60:     ky=iy/52774;
61:     iy=40692*(iy-ky*52774)-ky*3791;
62:
63:     rn=(float)(ix-iy)/2147483563.;
64:     rn-=(int)rn;

```

```

65:   if( rn<0.) rn++;
66:
67:   return rn;
68: }
69: /* ----- */
70: void sort(float x[],int n)
71: {
72:   int i,j,m;
73:   float xmin=10000000000.0;
74:   float w[100001];
75:
76:   for(i=1;i<=n;i++) w[i]=x[i];
77:   for(i=1;i<=n;i++){
78:     y[i]=xmin;
79:     for(j=1;j<=n-i+1;j++){
80:       if( y[i] > w[j] ){
81:         y[i]=w[j];
82:         m=j;
83:       }
84:     }

```

```

85:     for(j=m;j<=n-i;j++) w[j]=w[j+1];
86:   }
87: }

```

In Line 73, a large value is given to **xmin**.

In Lines 78 – 84, given **i**, we search the *i*th minimum value.

in Line 85, remove the *i*th minimum data and reconstruct the new vector.

Quick Sort: In the previous sort algorithm, the computational number of times is given by $n + (n - 1) + (n - 2) + \dots + 1 = n(n - 1)/2$, which is $O(n^2)$. That is, computational time is proportional to n^2 .

Computational time of the following sort algorithm, called the quick sort (クイック・ソート), is proportional to $n \log n$.

The following sort algorithm is extremely faster than the previous one in order of $n/\log n$.